# Quantum Error Correction: Dream or Nightmare

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### Back in 1982 Quantum Mechanical Hamiltonian Models of Turing Machines

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Quantum mechanical Hamiltonian models, which represent an aribtrary but finite number of steps of any Turing machine computation, are constructed here on a finite lattice of spin-1/2 systems. Different regions of the lattice correspond to different components of the Turing machine (plus recording system). Successive states of any machine computation are represented in the model by spin configuration states. Both time-independent and time-dependent Hamiltonian models are constructed here. The time-independent models do not dissipate

Time-independent Hamiltonian on a 2D lattice to execute a 1D quantum circuit (Lloyd & Terhal: Adiabatic and time-independent universal computing on a 2D lattice with simple 2-qubit interactions, New Journal of Physics 2016). Quantum Error Correction?

### Back in 1996...



After Peter Shor's factoring algorithm came out...

# Serge Haroche & Jean-Michel Raimond wrote in Physics Today



### QUANTUM COMPUTING: DREAM OR NIGHTMARE?

ples of quantum ig were laid out ars ago by comsts applying the on principle of schanics to comtion. Quantum ias recently betopic in physics, sognition that a dem can be pre-

Recent experiments have deepened our insight into the wonderfully counterintuitive quantum theory. But are they really harbingers of quantum computing? We doubt it.

Serge Haroche and Jean-Michel Raimond

two interacting qu trol" bit and a ' The control re changed, but its mines the evolutio get: If the co nothing happens t if it is 1, the targe a well-defined trar Quantum me mits additional a



# Their main points of criticism



- To do a computation with N operations and get a sensible answer, the error rate in each step should scale as 1/N or less. Such low error rates (10<sup>-10</sup> or less) are unphysical.
- Watchdog strategies or quantum error-correction is an experimenter's nightmare due to its complexity.
- Computing is different from creating coherent macroscopic quantum states, i.e. Bose-Einstein condensate (or superconducting state) as it involves information and manipulation.

### What is error correction?







J. von Neumann



From 1956





# What is error correction?



2D Ferromagnetic Ising model. Below critical temperature  $T_c$ : symmetry-breaking and stable magnetization.

Errors= spin flips

Ferromagnetic 2-spin interactions = 'parity checks'

Encode a qubit into a 2D Ising model?

$$|0\rangle = |\uparrow\uparrow\cdots\uparrow\rangle, |1\rangle = |\downarrow\downarrow\cdots\downarrow\rangle.$$

But a rotation  $e^{-iS_Z\pi/2}$  on a single spin can map  $|\uparrow\uparrow\cdots\uparrow\rangle + |\downarrow\downarrow\cdots\downarrow\rangle$  onto  $|\uparrow\uparrow\cdots\uparrow\rangle - |\downarrow\downarrow\cdots\downarrow\rangle$  $(S_Z|\uparrow\rangle = |\uparrow\rangle, S_Z|\downarrow\rangle = -|\downarrow\rangle).$ QUANTUM INFORMATION



### Models of Computation



Quantum

computing

(expected)

Universal

quantum

computing

Classical

(real)

computing







Noise Threshold: 0.6%- 1% error rate for each component e.g. DiVincenzo architecture for surface code using microwave resonators and transmon qubits



### The conundrum of small codes

Three-bit repetition code

$$|\overline{0}\rangle = |000\rangle, |\overline{1}\rangle = |111\rangle.$$

Parity checks are  $Z_1Z_2$  and  $Z_2Z_3$  measured non-destructively, e.g.



Using notation  $S_Z = Z$ ,  $S_X = X$ 

Single X errors detected and corrected. In quantum code we also measure also parity X-checks!



### The conundrum of small codes

Seven qubit code (Steane) encoding 1 qubit, able to correct a single error.





Parity check circuits

Nigg et al, Science (2014)

### The conundrum of small codes

Seven qubit code (Steane) encoding 1 qubit, able to correct a single error.





Parity check circuits

The means through which you get parity info. can also be the means through which you mess up your qubit!

# Why The Surface Code

Surface code for storing 1 logical qubit using d<sup>2</sup> physical qubits. Can correct d/2 errors (and more)

Smallest one: d=3 Surface-17 Below d=6



Qubits on vertices. Black squares=XXXX checks White squares=ZZZZ checks



Measure of encoding success? Get a encoded qubit with a longer lifetime  $\tau$  $(F(t) \approx e^{-t/\tau})$ . How fast are the encoded gates,  $t_{gate}$  (QEC slows things down!)? Improve  $\frac{t_{gate}}{\tau}$ !

# Logic

### 

All quantum power comes from the T gate

$$\mathrm{T}=\left(egin{array}{cc} 1 & 0 \\ 0 & e^{i\pi/4} \end{array}
ight)$$

When implementing universal QC with T gates one needs to process error information online, without running a backlog.



FIG. 6 Using the ancilla  $T|+\rangle$  in the dashed box, one can realize the T gate by doing a corrective operation SX.

For 2D stabilizer codes you cannot do the T gate via a constant-depth fault-tolerant circuit (Bravyi, Koenig 2013). Thus lots of overhead via 'magic-state-distillation technique'. In a 3D color code you can do a T gate without extra qubit overhead (smallest example [[15,1,3]]) but threshold is likely much worse than 1%.

Fowler et al., Phys. Rev. A 86, 032324 (2012)

#### Some Numbers/Estimates

Is space-time volume to factor N=2000 bit number realistic?

- Factoring a number with N=2000 bits needs 40 N<sup>3</sup> =O(10<sup>11</sup>) Toffolis (modular exponentiation) and about 2N=4000 logical qubits
- Each logical qubit (surface code) uses p=14,500 physical qubits (assume physical error rate=threshold/10), so 58 Mqbits
- Ancilla Factory. Each Toffoli needs 7 encoded T ancillas, so O(10<sup>12</sup>) encoded ancillas. Generating and purifying one ancilla takes 800,000 physical qubits (and 500 surface code cycles).

# Qubit into a microwave mode

Lots of space in a harmonic oscillator...

$$H = \hbar\omega(a^{\dagger}a + 1/2)$$

What states offer 'protection', form a code?







### **Displacement Sensor**

Assume weak time-dependent unknown force F(t) on oscillator so Hamiltonian  $H(t) = \hbar \omega (a^{\dagger}a + 1/2) - \hat{q}F(t)$ .

For example LC oscillator

$$H(t) = \hbar\omega(a^{\dagger}a + 1/2) + g V(t)(a + a^{\dagger}), \hat{q} = q = \frac{1}{\sqrt{2}}(a + a^{\dagger})$$
  

$$G_{in}$$
  

$$V(t) e.g.$$
  

$$= V_0 \cos(\omega' t + \varphi) = = = = = = LC \text{ oscillator } \omega = \sqrt{LC}$$

What are the limits in determining the displacement caused by V(t)?

### **Fundamental Limit?**

From Rev. Mod. Phys. (1980)

### On the measurement of a weak classical force coupled to a quantum-mechanical oscillator. I. Issues of principle\*

Carlton M. Caves, Kip S. Thorne, Ronald W. P. Drever,<sup>†</sup> Vernon D. Sandberg,<sup>‡</sup> and Mark Zimmermann<sup>§</sup>

#### B. Uncertainty principle and ways to measure the oscillator

In classical theory it is possible to measure the oscillator's complex amplitude  $X = X_1 + iX_2$  with complete precision. Not so in quantum theory. Equations (2.1) and (2.5) imply that  $\hat{X}_1$  and  $\hat{X}_2$  do not commute:

$$[\hat{X}_1, \hat{X}_2] = i\hbar/m\omega . \tag{2.8}$$

Therefore the variances of  $\hat{X}_1$  and  $\hat{X}_2$  in any oscillator state must satisfy

$$\Delta X_1 \Delta X_2 \ge \frac{1}{2} |\langle [\hat{X}_1, \hat{X}_2] \rangle| = \hbar/2m\omega , \qquad (2.9a)$$

which is the complex-amplitude analog of the Heisenberg uncertainty principle for position and momentum:

$$\Delta x \Delta p \ge \frac{1}{2}\hbar. \tag{2.9b}$$

But why measure p and q? We want to measure 2 parameters. Fundamental quantum limit is subtle. Terhal, Duivenvoorden, Single-mode Displacement Sensor, arXiv:org:1603.02242, submitted to PRL

## **Displacement Sensor**

Grid state  $|\psi_{grid}\rangle$  is a highly sensitive displacement sensor  $p \approx 0 \mod \sqrt{2\pi}$ ,  $q \approx 0 \mod \sqrt{2\pi}$ .



 $\overline{n} \approx \frac{1}{2\Delta^2}$  $\sigma$  of Gaussian envelope  $\sim \frac{1}{\Delta}$ and  $\sigma$  of individual peaks  $\sim \Delta$ 

Maximum strength of displacement on vacuum input  $\overline{n} \leq \pi/2$ 

Grid states introduced by Gottesman, Preskill, Kitaev in 2001 for quantum error correction.

### How well can one do?

Using Quantum Cramer-Rao Bound one can find for estimates  $\tilde{u}$  and  $\tilde{v}$  (of the parameters u and v in displacement  $e^{-i u \hat{p} + i v \hat{q}}$ )

 $Var(\tilde{u}) + Var(\tilde{v}) \ge 2$  (for coherent/thermal/squeezed states)

 $Var(\tilde{u}) + Var(\tilde{v}) \rightarrow \frac{1}{(2\bar{n}+1)}$  for 2-mode squeezed (EPR) state, one mode undergoing displacement

**Our Result** 

 $Var(\tilde{u}) + Var(\tilde{v}) = O(\frac{1}{\sqrt{n}})$  for grid state with phase estimation 'parity' measurement, for small u, v.

Cats in cavities, e.g. Vlastakis et al., Science 2013, Ofek et al.: arXiv.org:1602.04768

### Sensor state in Circuit-QED Hardware

- High-Q micro-cavity, say,1 msec or more.
- High quality qubit, say,  $T_1, T_2 \approx O(10 100) \, \mu sec$
- Strong dispersive qubit-cavity coupling  $\chi Za^{\dagger}a$ (e.g.  $\frac{\chi}{2\pi} = 2.5MHz$ , cavity/qubit detuning 1 GHz, non-linearities O(1) kHz)
- Dispersive coupling allows for qubit-controlled cavity rotation ( $R(\theta Z) = \exp(-i\theta \ a^{\dagger}a \ Z)$ ) which can be directly used for qubit-controlled displacement.



- Controlled-rotations take  $T = \pi/\chi = 200$  nanosec.
- Use no more than 50 photons. Squeezing required.
- Initial schemes worked out in Terhal/Weigand, PRA 2016.





### Conclusion

Creation of Grid or GKP code states may be experimentally feasible. They can be useful for encoding a qubit into an oscillator as well as for displacement sensing.